

# Combinatorics, Lecture 11, 2022/8/28

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§ 1.

Markov's inequality. Let  $X \geq 0$  be random variable and  $t > 0$ . Then  $\Pr(X \geq t) \leq \frac{\mathbb{E}[X]}{t}$ .

Pf.  $\mathbb{E}[X] = \sum_{\omega} X(\omega) \cdot P(\{\omega\})$

$$\geq \sum_{\substack{\omega: X(\omega) \geq t}} t \cdot P(\{\omega\})$$
$$= t \cdot \Pr(X \geq t). \quad \square$$

Corollary. Let  $X_n \geq 0$  be an integer-value random variable. If  $\mathbb{E}[X_n] \rightarrow 0$  as  $n \rightarrow +\infty$ ,

then  $\Pr(X_n = 0) \rightarrow 1$  as  $n \rightarrow +\infty$

Lemma 1. Let  $p = p(n)$  be a function of  $n$  and let  $G$  be  $G(n, p)$ . Then

$$\overline{\text{P}} \left( \alpha(G) \leq \left\lceil \frac{2 \ln n}{p} \right\rceil \right) \rightarrow 1 \text{ as } n \rightarrow \infty.$$

Pf. Let  $t = \left\lceil \frac{2 \ln n}{p} \right\rceil$ . Let  $X_n$  be the number of independent set of size  $t+1$  in  $G = G(n, p)$ . Then we have

$$\begin{aligned} \mathbb{E}[X_n] &= \sum_{S \subseteq \binom{[n]}{t+1}} \mathbb{E} [1_{\{S \text{ is independent}\}}] \\ &= \binom{n}{t+1} \cdot (1-p)^{\binom{t+1}{2}} \leq \frac{n}{(t+1)!} e^{-p^{\frac{t+1}{2}}} \\ &= \frac{1}{(t+1)!} \left( n e^{-p^{\frac{t}{2}}} \right)^{t+1} \\ t = \left\lceil \frac{2 \ln n}{p} \right\rceil &\leq \frac{1}{(t+1)!} \rightarrow 0 \quad (n \rightarrow \infty) \end{aligned}$$

$$\left( \begin{array}{c} \downarrow \\ t \geq 2 \ln n \\ \rightarrow +\infty \end{array} \right)$$

Def. The chromatic number

of a graph  $G$ , denoted by  $\chi(G)$ , is the least integer  $k$  such that  $V(G)$  can be

partitioned into  $k$  independent sets.

Fact 1  $\forall G, \chi(G) \cdot \alpha(G) \geq |V(G)|$ .

Thm 1 (Erdős) For any fixed  $k \in \mathbb{N}^+$ , there

exists a graph  $G$  with  $\chi(G) \geq k$  and  $g(G) \geq k$ .

Here,  $g(G)$  denotes the girth of  $G$ , i.e. the length of the shortest cycle in  $G$ .

Pf. Consider a random graph  $G = G(n, p)$ ,

where  $p$  will be determined later and  $n \gg k$ .

Let  $t = \lceil \frac{2 \ln n}{p} \rceil$ . By Lemma 1,

$$\Pr(\alpha(G) \leq t) \rightarrow 1 \text{ as } n \rightarrow \infty.$$

Let  $Y_n$  be the total number of cycles of length less than  $k$  in  $G$ .

$$\text{Then } E[Y_n] = \sum_{i=3}^{k-1} \frac{n(n-1)\cdots(n-i+1)}{2^i} p^i$$

$$\leq \sum_{i=3}^{k-1} (np)^i \leq \frac{(np)^{k-1}}{np - 1}$$

By Markov's inequality

$$P(Y_n > \frac{n}{2}) \leq \frac{E[Y_n]}{n/2} \leq \frac{2(np)^{k-1}}{n(np-1)}$$

$$\text{Let } p = n^{-\frac{k-1}{k}}. \text{ So } np = n^{\frac{1}{k}}$$

$$\Rightarrow P(Y_n > \frac{n}{2}) \leq \frac{2(n-1)}{n(n^{1/k-1})} \rightarrow 0$$

as  $n \rightarrow +\infty$

Then when  $n$  is sufficiently large, there exist a graph  $G$  with  $\Delta(G) \leq \lceil \frac{2 \ln n}{p} \rceil \leq 3n^{\frac{k-1}{k}}$

$$\text{and } Y_n \leq \frac{n}{2}.$$

By deleting one vertex from each cycle of

length at most  $k-1$ , we can find an induced subgraph  $G^*$  of  $G$ , which has at least  $\frac{n}{2}$  vertices and has no cycle of length  $\leq k-1$ .

Then  $\chi(G^*) \geq k$  and

$$\alpha(G^*) \leq \alpha(G) \leq 3n^{\frac{k-1}{k} \ln n}.$$

By Fact 1

$$\Rightarrow \chi(G^*) \geq \frac{|V(G^*)|}{\alpha(G^*)} \geq \frac{n/2}{3n^{\frac{k-1}{k} \ln n}} = \frac{n^{\frac{1}{k}}}{6 \ln n}$$

$$\geq k \quad (\text{as } n \gg k)$$

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Def. An  $k$ -uniform hypergraph  $G = (V, E)$

( $k$ -graph, for short) if  $E \subseteq \binom{V}{k}$ .

Def. We say a  $k$ -graph  $G$  is  $r$ -colorable if its vertices can be colored by  $r$  colors so that no edge is monochromatic.

Def. Let  $m(k)$  be the minimum number of edges in a  $k$ -graph which is not 2-colorable.

Thm (Erdős)  $m(k) \geq 2^{k-1}$ , i.e. any  $k$ -graph with fewer than  $2^{k-1}$  edges is 2-colorable.

Pf. Consider a random 2-coloring on vertices. The probability that an edge is monochromatic

$$= \frac{1}{\binom{k}{2}} = \frac{1}{\frac{k(k-1)}{2}} = \frac{2}{k(k-1)}$$

By union bound,  $\Pr(\exists \text{ a monochromatic edge})$

$$\leq \sum_{x=1}^{\infty} \frac{1}{2^{k-1}} < 1. \quad \blacksquare$$

## §2. Lovász Local Lemma

Consider "bad events"  $A_1, \dots, A_n$ .

We want to avoid all, i.e.  $\Pr(\overline{A_1 \cap \dots \cap A_n}) > 0$ .

- { ① If all  $\sum_i p(A_i)$  are small, say  $\sum_i p(A_i) < 1$ ,  
 then we can avoid all.  
 ② If they are all independent,  
 then  $p(\bar{A}_1 \cap \bar{A}_2 \cap \dots \cap \bar{A}_n) = (1 - p(A_1)) \dots (1 - p(A_n))$   
 $\Rightarrow (\text{if } \sum_i p(A_i) < 1)$

The LLL deals with the case when each bad event is independent with most other events, but possibly dependent with a small number of events.

Def. An event  $A_0$  is independent from  $\{A_1, \dots, A_m\}$   
 if  $p(A_0 \cap B_1 \cap \dots \cap B_m) = p(A_0) p(B_1 \cap \dots \cap B_m)$ ,  
 where each  $B_i \in \{A_i, \bar{A}_i\}$ . 6

Lovasz Local Lemma (symmetric form) Let

$A_1, A_2, \dots, A_n$  be events with  $p(A_i) \leq p$  v i.

Suppose that each  $A_i$  is independent from a set of all other  $A_j$  except for at most  $d$  of them.

If  $e^{\beta(d+1)} \leq 1$ , then none of the events  $A_i$  occur with a positive probability,

Lovasz Local Lemma (The general form)

Let  $A_1, \dots, A_n$  be events. For  $i \in [n]$ , let

$N(i) \subseteq [n]$  be such that  $A_i$  is independent from  $\{A_j : j \notin N(i) \cup \{i\}\}$ . If

$x_1, \dots, x_n \in [0, 1)$  satisfy

$$\Pr(A_{ij}) \leq x_i \prod_{j \in N(i)} (1 - x_j) \quad \text{for } \forall i \in [n],$$

then none of the events  $A_i$  occur with probability at least  $\frac{m}{n} (1 - x_i)$ .  $\checkmark$

Pf (The general form  $\Rightarrow$  symmetric form)

Suppose we have  $e^{\beta \prod_{j \in N(i)} A_j} \leq 1$ .

Let  $x_i = \frac{1}{d+1} < 1$  for all  $i \in [n]$ .

Then  $x_i \prod_{j \in N(i) \setminus \{i\}} (1 - x_j) = \frac{1}{d+1} \prod_{j \in N(i) \setminus \{i\}} (1 - \frac{1}{d+1})$

$$\geq \frac{1}{d+1} \left(1 - \frac{1}{d+1}\right)^d > \frac{1}{e(d+1)} \geq p \geq p(x_i)$$

so the conditions of the general form are

satisfied  $\Rightarrow P(\bar{A}_1, \dots, \bar{A}_n) > 0$   $\square$

Pf of the general form. We will prove that

for any  $i \notin S \subseteq [n]$ ,

$$P(A_i \mid \bigcap_{j \in S} \bar{A}_j) \leq x_i. \quad (*)$$

Once (x) is proved, we have

$$P(\bar{A}_1 \cap \bar{A}_2 \cap \dots \cap \bar{A}_n) = P(\bar{A}_1) \cdot P(\bar{A}_2 | \bar{A}_1) \cdot P(\bar{A}_3 | \bar{A}_1 \cap \bar{A}_2) \cdot \dots \cdot P(\bar{A}_n | \bar{A}_1 \cap \bar{A}_2 \cap \dots \cap \bar{A}_{n-1})$$

$$\Rightarrow (\neg x_1) \wedge (\neg x_2) \wedge \dots \wedge (\neg x_n)$$

Now we prove (x) by induction on |S|.

Base case : when |S|=0. It is trivial to

$$\text{get } P(A_{\emptyset}) \leq x_i$$

Consider  $i \notin S$ . Let  $S_1 = S \cup N^{(i)}$

$$S_2 = S \setminus N^{(i)}.$$

We have

$$P(A_i | \bigcap_{j \in S} \bar{A}_j) = \frac{P(A_i \cap \bigcap_{j \in S_1} \bar{A}_j | \bigcap_{j \in S_2} \bar{A}_j)}{P(\bigcap_{j \in S_1} \bar{A}_j | \bigcap_{j \in S_2} \bar{A}_j)}$$

①

The numerator of the above formula ①

$$\leq P(A_i \mid \bigcap_{j \in S_2} \overline{A_j}) = P(A_i)$$

$$\leq x_i \prod_{j \in N^{(i)}} (1 - x_j)$$

②

Let  $S_i = \{j_1, \dots, j_r\} \subseteq N^{(i)}$ .

Then the denominator of ①

$$\text{Let } B = \bigcap_{j \in S_2} \overline{A_j}$$

$$= P(\overline{A_j} \mid B) \cdot P(\overline{A_{j_2}} \mid \overline{A_j}, nB) \cdots \\ P(\overline{A_{j_r}} \mid \overline{A_j}, n \cdots n \overline{A_{j_{r-1}}}, nB)$$

by induction  
 $\geq (1 - x_{j_1}) (1 - x_{j_2}) \cdots (1 - x_{j_r})$

$$\geq \prod_{j \in N^{(i)}} (1 - x_j)$$

③

By ①, ②, ③

$$\Pr(A_i \mid \bigcap_{j \in S} \overline{A_j}) \leq \frac{\chi_i \prod_{j \in N(i) \setminus S} (1 - \chi_j)}{\prod_{j \in N(i) \setminus S} (1 - \chi_j)} = \chi_i$$

thus prove (t)

Thm 3 (Spencer, 1977) If  $e\left(\binom{k}{2}\binom{n}{k-2} + 1\right) 2^{1-\binom{k}{2}} < 1$ ,  
then  $R(k, k) > n$ .

Pf. Random 2-edge-coloring on  $K_n$ .

For each  $R \in \binom{[n]}{k}$ , let  $E_R$  be the event

that  $R$  induces a mono-chromatic  $K_k$ .

$$\Rightarrow \Pr(E_R) = 2^{1-\binom{k}{2}} \leq p.$$

For each  $R$ , it is independent of all  $E_S$

if  $|S \cap R| \leq 1$ .

$\Rightarrow$  There are at most  $\binom{k}{r} \binom{n}{k-r}$  sets  $\binom{[n]}{k}$  satisfying  $|SNR| \geq 2$ .

$$\Rightarrow d = \binom{k}{r} \binom{n}{k-r}.$$

By LLL (symmetric form) - expected  $< 1$

$$\Rightarrow |P\left(\bigcap_{i=1}^n E_i\right)| > 0 \Rightarrow R(k, k) > n$$

Corollary.  $R(k, k) > \left(\frac{\sqrt{e}}{e} + o(1)\right) k^{2^{-k/2}}$ .

This still is the best lower bound for  $R(k, k)$ !

Theorem A  $k$ -graph is 2-colorable if every edge intersects at most  $d = e^{-1} 2^{k-1} - 1$  other edges.

Pf: For each edge  $f$ , let  $A_f$  be the event that  $f$  is monochromatic (for random 2-coloring) on vertices.

$$\Rightarrow |P(A_f)| = 2^{1-k}.$$

Each  $A_f$  is independent from all  $A_{f'}$

where  $f' \cap f = \emptyset$ .

By condition, at most  $d = e^{-1} 2^{k-1} - 1$  edges

$$\text{intersecting with } f, \text{ and } \exp(d+1) = e 2^{k-1} (e^{-1} 2^{k-1}) \leq 1.$$

By LLL, none of the events  $A_f$  occur with positive probability.

$\Rightarrow \exists$  2-coloring on vertices such that

No edge is monochromatic

$\Rightarrow$  this  $k$ -graph is 2-colorable.  $\square$