

§.

Def. A probability space is a pair  $(\Omega, P)$ , where  $\Omega$  is a finite set and  $P: 2^\Omega \rightarrow [0, 1]$  is a function satisfying the following conditions:

$$(1) \quad P(\emptyset) = 0, \quad (2) \quad P(\Omega) = 1,$$

$$(3) \quad P(A \cup B) = P(A) + P(B) \text{ for disjoint sets } A, B \subseteq \Omega.$$

Def. A subset  $A$  of  $\Omega$  is called an event.

$$\text{and } P(A) = \sum_{w \in A} P(\{w\})$$

A random variable - function  $X: \Omega \rightarrow \mathbb{R}$ .

Def. The expectation of a random variable  $X$

$$\text{is } \mathbb{E}[X] = \sum_{w \in \Omega} P(\{w\}) \cdot X(w).$$

Fact 1. (Union bound)  $\forall A, B \subseteq \Omega$ ,

$$P(A \cup B) \leq P(A) + P(B)$$

Fact 2. (The linearity of expectation). For any random variables  $X$  and  $Y$ , we have

$$\mathbb{E}[X+Y] = \mathbb{E}[X] + \mathbb{E}[Y].$$

## §1. Union bound

Def. Recall  $R(s, t)$  (Ramsey number)

Thm 1. Let  $n, s$  satisfy  $\binom{n}{s} 2^{1 - \binom{s}{2}} < 1$ .

Then  $R(s, s) > n$ .

Pf. We need to find a 2-edge-coloring of  $K_n$  such that it has NO monochromatic cliques  $K_s$ .

Consider a random 2-edge-coloring  $c$  of  $K_n$ ,

where each edge of  $K_n$  is colored by red or blue, each with probability  $\frac{1}{2}$  and independent of each other.

(This is the same as the following: Let  $\mathbb{I}$  be the family of all 2-edge-colorings of  $K_n$ , and let  $c \in \mathbb{I}$  be obtained by choosing uniformly at random.)

Let  $B$  be event that this coloring  $c$  has NO monochromatic  $K_s$ . We want to show  $P(B) > 0$ .

Consider the complement event  $A = \mathbb{I} \setminus B$ , i.e.

$A$  is the event that  $c$  has a monochromatic  $K_s$ .

For  $S \subseteq \binom{[n]}{s}$ , let  $A_S$  be the event that

$S$  forms a monochromatic  $K_5$  for  $C$ .

$$\Rightarrow A = \bigcup_{S \in \binom{[n]}{5}} A_S \text{ and } P(A_S) = 2^{1 - \binom{|S|}{2}}$$

By the Union bound,

$$P(A) = P\left(\bigcup_{S \in \binom{[n]}{5}} A_S\right) \leq \sum_{S \in \binom{[n]}{5}} P(A_S)$$

$$= \binom{n}{5} \cdot 2^{1 - \binom{|S|}{2}} < 1.$$

i.e.  $P(A) > 0$ .



Consequently,  $R(n, 5) > \frac{1}{e^5} \cdot 5^5 \approx 1$ .

pf. (Exercise). choose a proper  $n = n(s)$  so that  
the condition of Thm 1 satisfies.

Def. The random graph  $G(n, p)$  for some  $p \in (0, 1)$   
is a graph with vertex set  $[n]$ , where each of the  
potential  $\binom{n}{2}$  edges appears with probability  $p$ , independent  
of each other. (Often, we consider  $p = \frac{1}{2}$ ).

- Let  $A$  be an event / a graphic property that we  
are interested in. Let

$p(A) = p(G(n, \frac{1}{2}) \text{ satisfies the property } A)$

a function on  $n$ ,

$$= \frac{\# \text{ of } n\text{-vertex graphs on } [n] \text{ satisfying } A}{2^{\binom{n}{2}}}$$

Def. We say  $G(n, \frac{1}{2})$  almost surely satisfies

a property  $A$ , if  $\lim_{n \rightarrow \infty} p(A) = 1$

Thm<sup>2</sup>.  $G(n, \frac{1}{2})$  almost surely is non-bipartite.

Pf. Let  $A$  be the event that  $G(n, \frac{1}{2})$  is bipartite.  
So we need to show  $\lim_{n \rightarrow \infty} p(A) = 0$ .

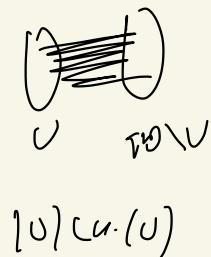
Fix a subset  $U \subseteq [n]$ , let  $A_U$  be the event that  $G(n, \frac{1}{2})$  is bipartite with bipartition  $(U, [n] \setminus U)$ .

$$A = \bigcup_{U \subseteq [n]} A_U$$

We have

$$p(A_U) = \frac{\# \text{ graphs on } [n] \text{ satisfying } A_U}{2^{\binom{n}{2}}}$$

$$= \frac{2^{(U)(n-U)}}{2^{\binom{n}{2}}}$$



$$\leq \frac{\frac{n^2}{4}}{\frac{2(n^2n)}{n}} = 2 - \frac{n^2}{4} + \frac{n}{2}$$

By Union bound,

$$0 \leq p(A) \leq \sum_{U \subseteq W} p(A_U) \leq 2^n \cdot 2 - \frac{n^2}{4} + \frac{n}{2}$$

$$= 2 - \frac{n^2}{4} + \frac{3}{2}n$$

$$\Rightarrow \lim_{n \rightarrow \infty} p(A) = 0.$$

Def. Given a probability space  $(\Omega, P)$ , events

$A_1, A_2, \dots, A_k$  are independent, if for any

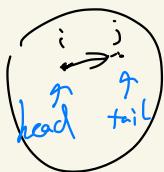
$I \subseteq [k]$ , we have  $p(\bigcap_{i \in I} A_i) = \prod_{i \in I} p(A_i)$ .

Def A tournament on  $n$  vertices is a

directed graph obtained from  $K_n$  by

assigning a direction to each edge.

An arc is an edge in a directed graph.



Def A tournament satisfies the property  $S_k$  if

for any subset  $S$  of size  $k$ , there exists

a vertex  $y \notin S$  such that  $y \rightarrow x$ .

for any  $x \in S$ ,



Thm3. For  $\forall k$ , if  $\boxed{\binom{n}{k} \left(1 - \frac{1}{2^k}\right)^{n-k} < 1}$ ,  
then there exists a tournament on  $n$  vertices satisfying  
the property  $S_k$ . T

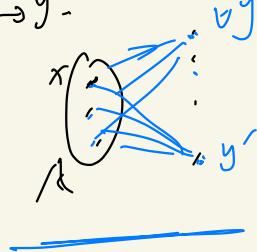
Pf. We consider a random tournament on  $[n]$ ,  
where for any pair  $\{i, j\}$ , the arc  $i \rightarrow j$  occurs with  
probability  $\frac{1}{2}$ , independent of each other.

Let  $B$  be the event that T does not satisfy

property  $S_k$ . For fixed  $A \subseteq \binom{[n]}{k}$ , let

$B_A$  be the event that the  $k$ -set  $A$  "fails",  
i.e.  $\forall y \in [n] \setminus A$ ,  $\exists x \in A$  with  $x \rightarrow y$ .

$$\Rightarrow B = \bigcup_{A \subseteq \binom{[n]}{k}} B_A$$



For  $y \in [n] \setminus A$ , let  $B_{A,y}$  be the event

that  $\exists x \in A$  with  $x \rightarrow y$ .

$$P(B_{A,y}) = 1 - \left(\frac{1}{2}\right)^k$$

Note that  $B_A = \prod_{y \in [n] \setminus A} B_{A,y}$  and events

$B_{A,y}$ 's are independent (as the  $k$  arcs between  $y$  and  $A$  are distinct for each  $y$ ).

$$\Rightarrow P(B_A) = \prod_{y \in T \setminus A} P(B_{A,y}) = \left(1 - \frac{1}{2^k}\right)^{n-k}.$$

$$\Rightarrow P(B) \leq \sum_{A \subseteq \binom{[n]}{k}} P(B_A) \leq \binom{n}{k} \cdot \left(1 - \frac{1}{2^k}\right)^{n-k} < 1. \quad \boxed{\square}$$

## §2. The linearity of expectation

Facts. For two random variables  $X, Y$ ,

$$E[X+Y] = E[X] + E[Y], \quad P(X \geq E[X]) > 0$$

$$\text{and } P(X \leq E[X]) > 0.$$

Def. A set  $A$  is sum-free, if for any  $x, y \in A$ ,  $x+y \notin A$ . (That is,  $x+y=z$  has NO solution in  $A$ .)

Examples { $\lceil \frac{n}{2} \rceil + 1, \lceil \frac{n}{2} \rceil + 2, \dots, n$ } is sum-free.  
 {all odd integers in  $[n]$ } is sum-free.

Exercise. The maximum size of a sum-free subset in  $[n]$  is  $\lceil \frac{n}{2} \rceil$ .

Thm 4. Let  $A$  be a finite set of non-zero integers.  
 There exists a sum-free subset  $B \subseteq A$  with  $|B| \geq \frac{|A|}{3}$ .

Pf: We choose a large prime  $p$  such that

$p > |a|$  for  $a \in A$ . Consider  $\mathbb{Z}_p = \{0, 1, \dots, p-1\}$

and  $\mathbb{Z}_p^* = \{1, 2, \dots, p-1\}$ . Note that there exists  
 a sum-free subset  $S$  under  $\mathbb{Z}_p \pmod{p}$ :

$$S = \left\{ \lceil \frac{p}{3} \rceil + 1, \lceil \frac{p}{3} \rceil + 2, \dots, \lceil \frac{2p}{3} \rceil \right\}.$$

Claim. For  $x \in \mathbb{Z}_p^*$ ,  $A_x = \{a \in A : ax \pmod{p} \in S\}$   
 is sum-free subset of  $A$ .

Pf: Suppose  $a, b, c \in A_x$  with  $a+b=c$ .

Then  $ax \pmod{p}$ ,  $bx \pmod{p}$ ,  $cx \pmod{p} \in S$ .

$$\text{But } ax + bx = cx \pmod{p},$$

a contradiction to that  $S$  is sum-free under  $\mathbb{Z}_p$ . □

Next, we turn to find an  $x \in \mathbb{Z}_p^*$  with  $|A_x| \geq |A|/3$ .

We will find such  $x$  by choosing  $x \in \mathbb{Z}_p^*$  uniformly  
 at random.

Compute  $\mathbb{E}[|A_x|]$ ,

where  $|A_x| = \sum_{a \in A} \mathbf{1}_{\{ax \bmod p \in S\}}$

$$\begin{aligned} \text{So } \mathbb{E}[|A_x|] &= \mathbb{E}\left[\sum_{a \in A} \mathbf{1}_{\{ax \bmod p \in S\}}\right] \\ &= \sum_{a \in A} \mathbb{E}\left[\mathbf{1}_{\{ax \bmod p \in S\}}\right] \\ &= \sum_{a \in A} P(ax \bmod p \in S). \end{aligned}$$

We observe that for fixed  $a \in A \implies (a \in \mathbb{Z}_p^*)$   
 $\{ax : x \in \mathbb{Z}_p^*\} = \mathbb{Z}_p^*$ .

$$\text{Then } P(ax \bmod p \in S) = \frac{|S|}{|\mathbb{Z}_p^*|} \geq \frac{1}{3}.$$

$$\implies \mathbb{E}[|A_x|] \geq \frac{1}{3}|A|.$$

As  $P(|A_x| \geq \mathbb{E}[|A_x|]) > 0$ , there exists some  
 $A_x$  with  $|A_x| \geq \mathbb{E}[|A_x|] \geq \frac{1}{3}|A|$ .  $\square$

Def. Let  $\alpha(G)$  be the maximum size of an  
 independent set in a graph  $G$ .

Thm 5. For any graph  $G$ ,  $\alpha(G) \geq \sum_{v \in V(G)} \frac{1}{d(v)+1}$   
 where  $d(v)$  denotes the degree of the vertex  $v \in V(G)$ .

Pf.: Let  $S_n$  be the family of all permutations

$\pi: [n] \rightarrow [n]$ . Let  $V^L(G) = [n]$ .

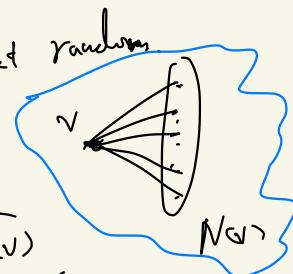
For a given permutation  $\pi$ , we say a vertex

$v \in V^L(G)$  is  $\pi$ -good if  $\pi(v)$  is the smallest value over the vertices  $\{v\} \cup N_G(v)$ .

Let  $\pi \in S_n$  be chosen uniformly at random

Then  $v \in V^L(G)$

$$P(v \text{ is } \pi\text{-good}) = \frac{1}{1 + d(v)}$$



Claim. All  $\pi$ -good vertices form an independent set.

Pf.: Suppose  $\exists$  two  $\pi$ -good vertices  $i, j$  with  $i, j \in V^L(G)$ .

Let  $\pi^{(i)} < \pi^{(j)}$ .  $\Rightarrow j$  is NOT  $\pi$ -good,  
a contradiction.  $\square$

Compute  $\left[ \mathbb{E} \left[ \left| \{ \text{all } \pi\text{-good vertices} \} \right| \right] \right]$

$$= \sum_{v \in V^L(G)} \mathbb{E} [1_{\{v \text{ is } \pi\text{-good}\}}]$$

$$= \sum_{v \in V^L(G)} P(v \text{ is } \pi\text{-good})$$

$$= \sum_{v \in V^L(G)} \frac{1}{1 + d(v)}$$

Then there exists some  $\pi \in S_n$  with

$$\alpha(G) \geq \left| \{ \text{all } \pi \text{-good vertices} \} \right| \geq E\left[ \left| \{ \text{good vertices} \} \right| \right] \\ \geq \sum_{v \in V(G)} \frac{1}{1 + d(v)}. \quad \boxed{\square}$$

Corollary. For any graph  $G$  on  $n$  vertices and  $m$  edges,

we have  $\alpha(G) \geq \frac{n^2}{2m+n}$ .

Pf. Exercise. □

Turán's Theorem (Approximate version) If  $G$  has  $n$  vertices and is  $K_r$ -free, then  $\alpha(G) \leq \frac{r-1}{2r} n$ .

2nd pf. Using previous result. □

### § 3. The deletion method

Earlier, we often define an approximate probability space and show that the random event occurs with positive probability.

Now we extend this idea, and consider situations where the random event does not always have the desired property, and may have very few "blemishes".

The point that we want to make here is that after deleting all blemishes, we will obtain the desired property

Recall. If  $\binom{n}{k} 2^{1-\frac{k}{n}} < 1$ , then  $R(k, k) > n$ .

$$\Rightarrow R(k, k) > \frac{1}{e^{\frac{1}{k}}} k 2^{\frac{k}{n}}$$

Thm 6. For any integer  $n$ , we have

$$R(k, k) > n - \binom{n}{k} 2^{1-\frac{k}{n}}$$

Pf. Consider a random 2-edge-coloring of  $K_n$ .

For  $A \in \binom{\text{edges}}{k}$ , let  $X_A$  be the indicator random variable of the event that  $A$  induces a monochromatic  $K_k$ . Then  $E[X_A] = 2^{1-\frac{k}{n}}$ .

Let  $X = \sum_{A \in \binom{\text{edges}}{k}} X_A$  be the number of

monochromatic  $K_k$ 's. Then

$$E[X] = \sum_{A \in \binom{\text{edges}}{k}} E[X_A] = \binom{n}{k} 2^{1-\frac{k}{n}}$$

So there exists a 2-edge-coloring of  $K_n$  whose number of monochromatic  $K_k$  is at most

$$\mathbb{E}[X] = \binom{n}{k} 2^{1 - \binom{k}{2}}$$

Next, we remove one vertex from each monochromatic  $K_k$ . This will delete at most

$X \leq \binom{n}{k} 2^{1 - \binom{k}{2}}$  vertices, and will destroy all monochromatic  $K_k$ . So it removes at least  $n - \binom{n}{k} 2^{1 - \binom{k}{2}}$  vertices.

with no monochromatic  $K_k$ .

$$\Rightarrow R(k, k) > n - \binom{n}{k} 2^{1 - \binom{k}{2}} \quad \text{✓}$$

Corollary .  $R(k, k) > \frac{1}{e}(1 + o(1)) k 2^{\frac{k}{2}}$

Pf.

Exercise:

✓

























